Linear and Network Optimization Exercise 8

Please return your solutions by Tuesday, June 3rd, 10:00 a.m., in the mailbox No. 5.

Problem 1 (7 points)

Formulate the simplex method for LPs with bounded variables as an algorithm, including the update of the reduced tableau.

Problem 2 (7 points)

Apply your algorithm to the LP $\min\{\underline{c} \underline{x} : A\underline{x} = \underline{b}, \underline{l} \leq \underline{x} \leq \underline{u}\}$ with

$$A = \begin{pmatrix} 1 & 1 & 0 & 0 & 0 \\ -1 & 0 & 1 & 1 & 0 \\ 0 & -1 & -1 & 0 & 1 \end{pmatrix}, \quad \underline{b} = \begin{pmatrix} 2 \\ 1 \\ -1 \end{pmatrix}$$
$$\underline{c} = (1, 1, 2, 3, 1), \qquad \underline{u} = (3, 3, 4, 2, 1)^T, \qquad \underline{l} = (0, 0, 0, 0, 0)^T$$

For the starting solution you can use

$$B = \{1, 3, 5\}, \quad L = \{2\}, \quad U = \{4\}.$$

Problem 3 (6 points)

Consider the Klee-Minty example of an LP where the simplex method goes through all of the 2^n basic feasible solutions in the worst case (with $0 < \varepsilon < \frac{1}{2}$):

(a) Show that the Klee-Minty problem is equivalent to the following problem, where $\theta := \frac{1}{\epsilon}$:

$$\max \sum_{j=1}^{n} y_{j}$$

s.t.
$$y_{j} + 2 \sum_{k=1}^{j-1} y_{k} \leq \theta^{j-1} \quad \forall j = 2, \dots, n$$
$$y_{j} \geq 0 \quad \forall j = 1, \dots, n$$

Hint: Use the transformation $y_1 := x_1$ and $y_j := \frac{(x_j - \varepsilon x_{j-1})}{\varepsilon^{j-1}}$ for $j = 2, \ldots, n$.

- (b) Graph the two problems for n = 2 and $\varepsilon = \frac{1}{3}$.
- (c) Apply the simplex method to the reformulation of subproblem (a) and illustrate the path of the algorithm through the basic feasible solutions in the graphs of subproblem (b).