

# Efficient points for Weber problems with asymmetric mixed distances

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- Majority principle for classical Weber problem

- Majority principle for (WAM)

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- Convex hull

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## Multicriteria perspective

- Multicriteria optimization problem

- Efficient points for (MOP)

## Algorithmic realization

- Algorithm to determine sets of efficient points

- Examples for strictly efficient sets



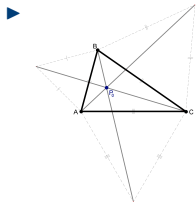
# Historical roots

- ▶ Pierre de Fermat (1601-1665):  
"... given three points in the plane, find a fourth point such that the sum of its distances to the three given points is a minimum ..." [Kuhn '67]



# Historical roots

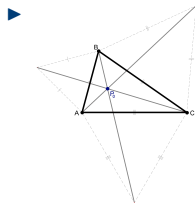
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- ▶ Battista Cavalieri (1598-1647): "Exerciones Geometricae"

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# Weber problem with asymmetric mixed distances

Mathematical model of the classical Weber problem :

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N d(x, a_n),$$

with

- ▶ existing facilities  $a_n \in A \subset \mathbb{R}^2$ ,  
 $n \in \mathcal{N} := \{1, \dots, N\}$
- ▶ uniform weights and
- ▶ a uniform symmetric distance measure  $d(\cdot, \cdot) : \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$  induced by a norm



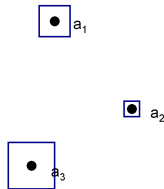
# Weber problem with asymmetric mixed distances

Mathematical model of the classical Weber problem :

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n d(x, a_n),$$

with

- ▶ existing facilities  $a_n \in A \subset \mathbb{R}^2$ ,  $n \in \mathcal{N} := \{1, \dots, N\}$
- ▶ **positive weights**  $w_n > 0, \forall n \in \mathcal{N}$
- ▶ a uniform symmetric distance measure  $d(\cdot, \cdot) : \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$  induced by a norm



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$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n d_n(x, a_n),$$

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 $n \in \mathcal{N} := \{1, \dots, N\}$
- ▶ positive weights  $w_n > 0, \forall n \in \mathcal{N}$
- ▶ symmetric distance measures  
 $d_n(\cdot, \cdot) : \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$  induced by  
not necessarily uniform norms





# Weber problem with asymmetric mixed distances

Mathematical model of the Weber problem with mixed asymmetric distances:

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n \gamma_n(x, a_n),$$

with

- ▶ existing facilities  $a_n \in A \subset \mathbb{R}^2$ ,  
 $n \in \mathcal{N} := \{1, \dots, N\}$
- ▶ positive weights  $w_n > 0, \forall n \in \mathcal{N}$
- ▶ asymmetric distance measures induced by (polyhedral) gauges  
 $\gamma_n(x, a) := \inf \{ \lambda > 0 : \frac{x}{\lambda} \in B_n + a \}$ ,  
with

- ▶  $B_n$  convex, closed and bounded
- ▶  $0 \in \text{int}(B_n)$
- ▶  $\gamma_n(x, y) \geq 0$
- ▶  $\gamma_n(x, y) = 0 \Leftrightarrow x = y$
- ▶  $\gamma_n(x, y) \leq \gamma_n(x, z) + \gamma_n(z, y)$



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# Majority principle for the classical Weber problem

## Theorem (Witzgall '64)

*For existing facilities  $a_n \in \mathbb{R}^2$ , weights  $w_n > 0$  for  $n \in \mathcal{N}$  and the classical Weber problem (with uniform symmetric distance measures)*

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n d(x, a_n)$$

*an optimal solution can be obtained in the existing facility  $a_k$  if*

$$w_k \geq \sum_{\substack{j \in \mathcal{N} \\ j \neq k}} w_j.$$

# Majority principle for (WAM)

## Theorem

An existing facility  $a_k$ ,  $k \in \mathcal{N}$  is an optimal solution for (WAM) if

a) all gauges  $\gamma_n$  are symmetrical and  $w_k B_k \supseteq \sum_{n \neq k} w_n B_n$  or

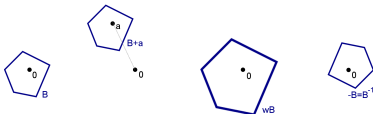
b)  $w_k B_k \supseteq \sum_{n \neq k} -w_n B_n =: \sum_{n \neq k} w_n B_n^{-1}$ .

Notations and conventions, with  $x, a \in \mathbb{R}^2$ ,  $w \in \mathbb{R}^+$  and  $B \subset \mathbb{R}^2$ :

►  $x \in B \Leftrightarrow x + a \in (B + a)$

►  $x \in B \Leftrightarrow wx \in wB$

►  $x \in B \Leftrightarrow -x \in B^{-1}$



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b)  $w_k B_k \supseteq \sum_{n \neq k} -w_n B_n =: \sum_{n \neq k} w_n B_n^{-1}$ .

Proof (sketch).

a)  $\blacktriangleright w_k B_k \supseteq \sum_{n \neq k} w_n B_n \Rightarrow w_k \gamma_k(x, a) \geq \sum_{n \neq k} w_n \gamma_n(x, a) \quad (1)$

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$$\begin{aligned} &\blacktriangleright \sum_{n \in \mathcal{N}} w_n \gamma_n(x, a_n) - \sum_{n \in \mathcal{N}} w_n \gamma_n(a_k, a_n) \\ &= \sum_{n \neq k} w_n (\gamma_n(x, a_n) - \gamma_n(a_k, a_n)) + w_k \gamma_k(x, a_k) \\ &\stackrel{(1)}{\geq} \sum_{n \neq k} w_n (\gamma_n(x, a_n) + \gamma_n(x, a_k) - \gamma_n(a_k, a_n)) \\ &= \sum_{n \neq k} w_n (\gamma_n(x, a_n) + \gamma_n(a_k, x) - \gamma_n(a_k, a_n)) \geq 0 \end{aligned}$$



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$$\text{b) } \blacktriangleright w_k B_k \supseteq \sum_{n \neq k} w_n B_n^{-1} \Rightarrow w_k \gamma_k(x, a) \geq \sum_{n \neq k} w_n \gamma_n(a, x) \quad (2)$$



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$$\begin{aligned} & \sum_{n \in \mathcal{N}} w_n \gamma_n(x, a_n) - \sum_{n \in \mathcal{N}} w_n \gamma_n(a_k, a_n) \\ &= \sum_{n \neq k} w_n (\gamma_n(x, a_n) - \gamma_n(a_k, a_n)) + w_k \gamma_k(x, a_k) \\ &\stackrel{(2)}{\geq} \sum_{n \neq k} w_n (\gamma_n(x, a_n) + \gamma_n(a_k, x) - \gamma_n(a_k, a_n)) \geq 0 \end{aligned}$$

## Theorem (Wendell '71)

Let  $A \subset \mathbb{R}^2$  be a finite set. Then for every  $x_1 \in \mathbb{R}^2$  and  $x_2 \in \text{conv}(A)$  exists with

$$\|x_1 - a_n\|_p \geq \|x_2 - a_n\|_p \quad \forall a_n \in A$$

for every  $1 \leq p \leq \infty$ .

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$$\|x_1 - a_n\|_p \geq \|x_2 - a_n\|_p \quad \forall a_n \in A$$

for every  $1 \leq p \leq \infty$ .

⇒ At least one optimal solution of the classical Weber problem (with uniform symmetric distance measures)

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n d(x, a_n)$$

lies within the convex hull of the existing facilities.

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## Theorem (Juel '83)

If  $x^*$  is an optimal solution of

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n d(x, a_n),$$

where  $d : \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$  is a distance measure induced by an  $\ell_p$  norm with  $1 < p < \infty$  and  $A := \{a_1, \dots, a_N\}$ , then  $x^* \in \text{conv}(A)$ .

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# Convex hull properties

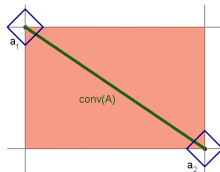
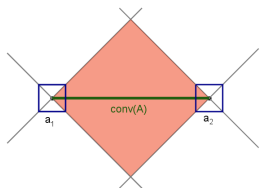
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Counterexample for  $p = \infty$  and  $p = 1$  ( $w_1 = w_2$  in both cases):



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# Counterexample for (WAM)

- ▶ For problem (WAM)

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n \gamma_n(x, a_n)$$

no analog property using  $\text{conv}(A)$  can be obtained.

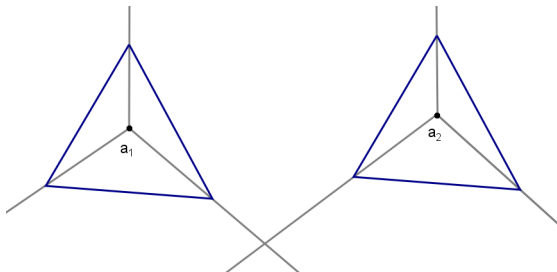
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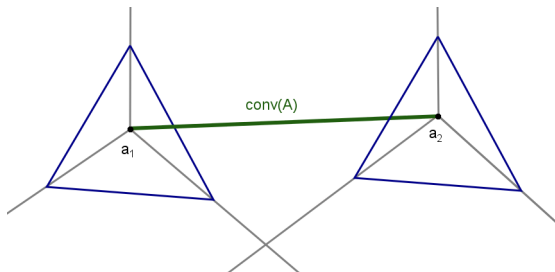
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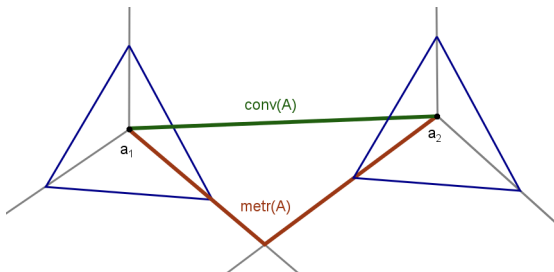
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- ▶ Example with  $w_1 = w_2$  and identical asymmetric distance measures:



## Definition

The metric hull of the set  $A$  of points  $a_n$ ,  $n \in \mathcal{N}$  is defined as

$$\text{metr}(A) = \{x \in \mathbb{R}^2 : \forall y \in \mathbb{R}^2, x \neq y, \\ \exists a_n \in A \text{ with } \gamma_n(x, a_n) < \gamma_n(y, a_n)\}$$

with respect to the gauges  $\gamma_n$ ,  $n \in \mathcal{N}$ .

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with respect to the gauges  $\gamma_n$ ,  $n \in \mathcal{N}$ .

## Basic properties: (Durier '85)

- ▶  $A \subset \text{metr}(A)$
- ▶  $\exists x^* \in \text{metr}(A)$  such that  $x^*$  is optimal for (WAM)

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# Convex and metric hull

## Further properties: (Durier '86)

Let  $X$  be a normed space and let the distance measures be induced by uniform  $\ell_p$  norms  $1 < p < \infty$ , then

$$\text{metr}(A) \subseteq \text{conv}(A).$$



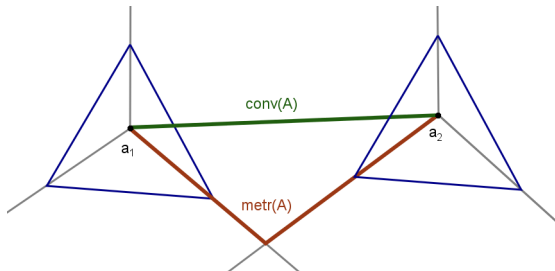
# Convex and metric hull

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This does not hold for (WAM):



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# Multicriteria optimization problem

Considering the multicriteria optimization problem

$$\min_{x \in \mathbb{R}^2} \begin{pmatrix} \gamma_1(x, a_1) \\ \vdots \\ \gamma_N(x, a_N) \end{pmatrix} \quad (\text{MOP})$$

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Considering the multicriteria optimization problem

$$\min_{x \in \mathbb{R}^2} \begin{pmatrix} \gamma_1(x, a_1) \\ \vdots \\ \gamma_N(x, a_N) \end{pmatrix} \quad (\text{MOP})$$

the problem (WAM)

$$\min_{x \in \mathbb{R}^2} \sum_{n=1}^N w_n \gamma_n(x, a_n)$$

can be interpreted as weighted sum scalarization with weights according to the customer demand.

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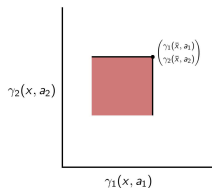
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# Efficient points for (MOP)

## Concepts of efficiency in multicriteria optimization:

- ▶  $\bar{x}$  is weakly efficient  $\Leftrightarrow \nexists x \in \mathbb{R}^2 : \gamma_n(x, a_n) < \gamma_n(\bar{x}, a_n)$   
 $\forall n \in \mathcal{N}$

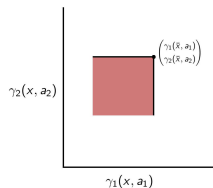


weakly efficient

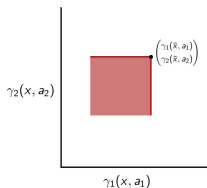
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 $\exists n \in \mathcal{N} : \gamma_n(x, a_n) < \gamma_n(\bar{x}, a_n)$



weakly efficient

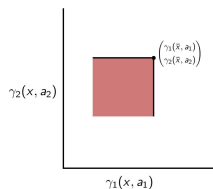


efficient

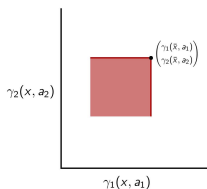
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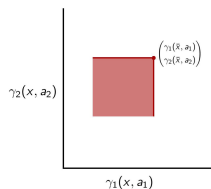
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- ▶  $\bar{x}$  is efficient  $\Leftrightarrow \nexists x \in \mathbb{R}^2 : \gamma_n(x, a_n) \leq \gamma_n(\bar{x}, a_n) \forall n \in \mathcal{N}$  and  
 $\exists n \in \mathcal{N} : \gamma_n(x, a_n) < \gamma_n(\bar{x}, a_n)$
- ▶  $\bar{x}$  is strictly efficient  $\Leftrightarrow \nexists x \in \mathbb{R}^2, x \neq \bar{x} :$   
 $\gamma_n(x, a_n) \leq \gamma_n(\bar{x}, a_n) \forall n \in \mathcal{N}$



weakly efficient



efficient



strictly efficient



# Boundedness of $X_{w\text{-eff}}$

## Theorem

*The set of weakly efficient points  $X_{w\text{-eff}}$  for (WAM) is bounded.*

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# Boundedness of $X_{W\text{-eff}}$

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*The set of weakly efficient points  $X_{W\text{-eff}}$  for (WAM) is bounded.*

## Proof (sketch).

- ▶ level sets  $L_{\gamma_n}(\lambda) := \{x \in \mathbb{R}^2 : \gamma_n(x, a_n) \in \lambda B_n\} \rightarrow$  convex and bounded

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- ▶ let  $\bar{L}_{\gamma_n} := \min\{L_{\gamma_n}(\lambda) : \gamma_n(a, a_n) \in \lambda B_n, \forall a \in A\}$

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- ▶ let  $\bar{L}_{\gamma_n} := \min\{L_{\gamma_n}(\lambda) : \gamma_n(a, a_n) \in \lambda B_n, \forall a \in A\}$
- ▶ define  $\bar{L} := \bigcup_{n \in \mathcal{N}} \bar{L}_{\gamma_n}$ , then  
 $\forall x \notin \bar{L}$  and  $\forall \bar{x} \in \bar{L} : \gamma_n(x, a_n) > \gamma_n(\bar{x}, a_n)$

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## Theorem

*The set of weakly efficient points  $X_{W-eff}$  for (WAM) is bounded.*

## Proof (sketch).

- ▶ level sets  $L_{\gamma_n}(\lambda) := \{x \in \mathbb{R}^2 : \gamma_n(x, a_n) \in \lambda B_n\} \rightarrow$  convex and bounded
- ▶ let  $\bar{L}_{\gamma_n} := \min\{L_{\gamma_n}(\lambda) : \gamma_n(a, a_n) \in \lambda B_n, \forall a \in A\}$
- ▶ define  $\bar{L} := \bigcup_{n \in \mathcal{N}} \bar{L}_{\gamma_n}$ , then  
 $\forall x \notin \bar{L}$  and  $\forall \bar{x} \in \bar{L} : \gamma_n(x, a_n) > \gamma_n(\bar{x}, a_n)$
- ▶  $\Rightarrow X_{W-eff} \subseteq \bar{L}$

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# Algorithm to determine sets of efficient points

Input

Step 1: Grid points

Step 2: Efficient grid points

Step 3: Efficient grid lines

Step 4: Efficient grid cells

Output

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# Algorithm to determine sets of efficient points

## Input

Existing facilities  $a_n \in A \subset \mathbb{R}^2$ ,  $n \in \mathcal{N}$  and associated distance measures  $\gamma_n$  induced by polyhedral gauges.

## Step 1: Grid points

## Step 2: Efficient grid points

## Step 3: Efficient grid lines

## Step 4: Efficient grid cells

## Output



# Algorithm to determine sets of efficient points

## Input

## Step 1: Grid points

Determine the set of grid points  $\mathcal{P}_G$  of the construction grid  $G$  based on the fundamental directions  $v_{nj}$  of the polyhedral gauges  $\gamma_n$ .

## Step 2: Efficient grid points

## Step 3: Efficient grid lines

## Step 4: Efficient grid cells

## Output





# Algorithm to determine sets of efficient points

Input

Step 1: Grid points

Step 2: Efficient grid points

Test all  $x \in \mathcal{P}_G$  if they are (weakly/strictly) efficient.

Set  $x \in M_P^{(w/s)eff} \Leftrightarrow x \in \mathcal{P}_G \wedge x \in X_{(w/s)-eff}$ .

Step 3: Efficient grid lines

Step 4: Efficient grid cells

Output

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# Algorithm to determine sets of efficient points

Input

Step 1: Grid points

Step 2: Efficient grid points

Step 3: Efficient grid lines

For all  $x \neq y \in M_P^{(w/s)eff}$  determine  $t_{xy} = \frac{1}{2}(x + y)$ .

If  $t \in X_{(w/s)-eff} \wedge \exists n, j : x, y, t \in v_{nj}$  set  $\overline{xy} \in M_L^{(w/s)eff}$ .

Step 4: Efficient grid cells

Output

# Algorithm to determine sets of efficient points

Input

Step 1: Grid points

Step 2: Efficient grid points

Step 3: Efficient grid lines

Step 4: Efficient grid cells

If  $\overline{x_i x_j} \in M_L^{(w/s)eff}$  yields for all line segments  $\overline{x_i x_j}$  that are bounding a construction grid cell  $D_\pi$ , then set  $D_\pi \in M_C^{(w/s)eff}$ .

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# Algorithm to determine sets of efficient points

Input

Step 1: Grid points

Step 2: Efficient grid points

Step 3: Efficient grid lines

Step 4: Efficient grid cells

Output

Set of (weakly/strictly) efficient points

$$X_{(w/s)\text{-eff}} = M_P^{(w/s)\text{eff}} \cup M_L^{(w/s)\text{eff}} \cup M_C^{(w/s)\text{eff}} .$$

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# Tests for (weakly/strictly) efficiency



definition

$$\begin{array}{|l} \bar{x} \in X_{w\text{-eff}} \\ \nexists x \in \mathbb{R}^2 : \gamma_n(x, a_n) < \gamma_n(\bar{x}, a_n) \quad \forall n \in \mathcal{N} \end{array}$$

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# Tests for (weakly/strictly) efficiency



definition	$\bar{x} \in X_{w\text{-eff}}$
	$\nexists x \in \mathbb{R}^2 : \gamma_n(x, a_n) < \gamma_n(\bar{x}, a_n) \quad \forall n \in \mathcal{N}$
hull of $\nabla\gamma_n(\bar{x}, a_n)$	$\bar{x} \in \text{int}(\text{conv}\{\bar{x} + \nabla\gamma_n(\bar{x}, a_n), n \in \mathcal{N}\})$

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# Tests for (weakly/strictly) efficiency



definition	$\bar{x} \in X_{w-eff}$ $\nexists x \in \mathbb{R}^2 : \gamma_n(x, a_n) < \gamma_n(\bar{x}, a_n) \quad \forall n \in \mathcal{N}$
hull of $\nabla\gamma_n(\bar{x}, a_n)$	$\bar{x} \in \text{int}(\text{conv}\{\bar{x} + \nabla\gamma_n(\bar{x}, a_n), n \in \mathcal{N}\})$
angles of $\nabla\gamma_n(\bar{x}, a_n)$	$\max\{\alpha_i + \alpha_j\} \geq 180^\circ$

where  $\alpha_i, \alpha_j$  are the directed angles between two gradients  $\nabla\gamma_{n_1}(\bar{x}, a_{n_1}), \nabla\gamma_{n_2}(\bar{x}, a_{n_2})$  and  $\nabla\gamma_{n_2}(\bar{x}, a_{n_2}), \nabla\gamma_{n_3}(\bar{x}, a_{n_3})$   $\forall n_1, n_2, n_3 \in \mathcal{N}$ , respectively.

# Tests for (weakly/strictly) efficiency



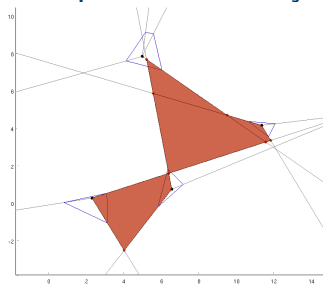
definition	$\bar{x} \in X_{w-eff}$ $\nexists x \in \mathbb{R}^2 : \gamma_n(x, a_n) < \gamma_n(\bar{x}, a_n) \forall n \in \mathcal{N}$
hull of $\nabla\gamma_n(\bar{x}, a_n)$	$\bar{x} \in \text{int}(\text{conv}\{\bar{x} + \nabla\gamma_n(\bar{x}, a_n), n \in \mathcal{N}\})$
angles of $\nabla\gamma_n(\bar{x}, a_n)$	$\max\{\alpha_i + \alpha_j\} \geq 180^\circ$

where  $\alpha_i, \alpha_j$  are the directed angles between two gradients  $\nabla\gamma_{n_1}(\bar{x}, a_{n_1}), \nabla\gamma_{n_2}(\bar{x}, a_{n_2})$  and  $\nabla\gamma_{n_2}(\bar{x}, a_{n_2}), \nabla\gamma_{n_3}(\bar{x}, a_{n_3})$   $\forall n_1, n_2, n_3 \in \mathcal{N}$ , respectively.

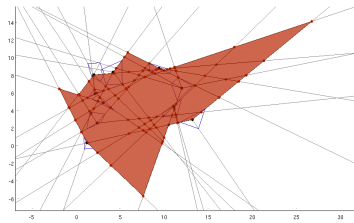
- ▶ analog conditions can be obtained for strictly efficient and efficient points



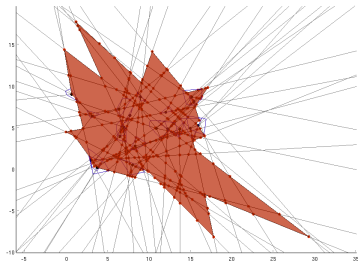
# Examples for strictly efficient sets



$N = 4$



$N = 10$



$N = 20$

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# Thank you for your attention!

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