

Standort-Optimierung

Handout 4

Sommersemester 2010



Bergische Universität Wuppertal

Fachbereich C – Angewandte Mathematik / Optimierung und Approximation
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Algorithm 5.3.: LP-Based Algorithm for $m/P/\bullet/l_1/\max$

Input: Existing facilities $a_1, \dots, a_n \in \mathbb{R}^2$; number m of new facilities sought; positive weights $w_{1,ij}$, $i = 1, \dots, m$, $j = 1, \dots, n$ and $w_{2,ir}$, $i = 1, \dots, m-1$, $r = i+1, \dots, m$.

Step 1: Find an optimal solution \tilde{x} of the corresponding problem of type $m/P/\bullet/l_\infty/\max$ with existing facility locations at $T^{-1}(a_1), \dots, T^{-1}(a_n)$.

Output: Optimal solution $x^* := T(\tilde{x})$ of $m/P/\bullet/l_1/\max$.

Algorithm 5.2: LP-Based Algorithm for $m/P/\bullet/l_\infty/\max$

Input: Existing facilities $a_1, \dots, a_n \in \mathbb{R}^2$; number m of new facilities sought; positive weights $w_{1,ij}$, $i = 1, \dots, m$, $j = 1, \dots, n$ and $w_{2,ir}$, $i = 1, \dots, m-1$, $r = i+1, \dots, m$.

Step 1: For $k = 1, 2$ do: Solve the LP

$$\begin{array}{ll}\min & z_k \\ \text{s.t.} & \left. \begin{array}{l} -x_{ik} + \frac{1}{w_{1,ij}}z_k \geq -a_{jk} \\ x_{ik} + \frac{1}{w_{1,ij}}z_k \geq a_{jk} \end{array} \right\} \forall i \in \{1, \dots, m\}, \\ & \left. \begin{array}{l} -x_{ik} + x_{rk} + \frac{1}{w_{2,ir}}z_k \geq 0 \\ x_{ik} - x_{rk} + \frac{1}{w_{2,ir}}z_k \geq 0 \end{array} \right\} \forall i, r \in \{1, \dots, m\}, \\ & \quad i < r \end{array}$$

and determine the optimal solution $x_{1k}^*, \dots, x_{mk}^*, z_k^*$.

Step 2: Set $x_i^* := (x_{i1}^*, x_{i2}^*)$, $i = 1, \dots, m$, $x^* := (x_1^*, \dots, x_m^*)$, and $z^* := \max\{z_1^*, z_2^*\}$.

Output: Optimal solution x^* of $m/P/\bullet/l_\infty/\max$ with objective value z^* .